

# Esquemas de actualización que preservan características dinámicas

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22 de abril de 2016

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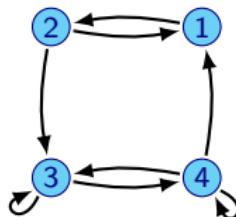
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# Motivation

## Boolean Networks

- A finite set  $V$  of  $n$  elements and  $n$  states variables  $x_v \in \{0, 1\}$ ,  $v \in V$
- A global activation function  $F = (f_v)_{v \in V} : \{0, 1\}^n \rightarrow \{0, 1\}^n$ 
  - Composed by local activation functions  $f_v : \{0, 1\}^n \rightarrow \{0, 1\}$
- Schedule



$$\begin{aligned} f_1(x) &= x_2 \wedge x_4 \\ f_2(x) &= x_1 \\ f_3(x) &= x_2 \vee x_3 \\ f_4(x) &= x_3 \wedge x_4 \end{aligned}$$

$$\begin{aligned} s(i) &= 1 && \iff s = \{1, 2, 3, 4\} \\ s(i) &= i && \iff s = \{1\} \{2\} \{3\} \{4\} \\ s(i) &= \begin{cases} 1 & \text{if } i > 2 \\ 2 & \text{if } i \leq 2 \end{cases} && \iff s = \{3, 4\} \{1, 2\} \end{aligned}$$

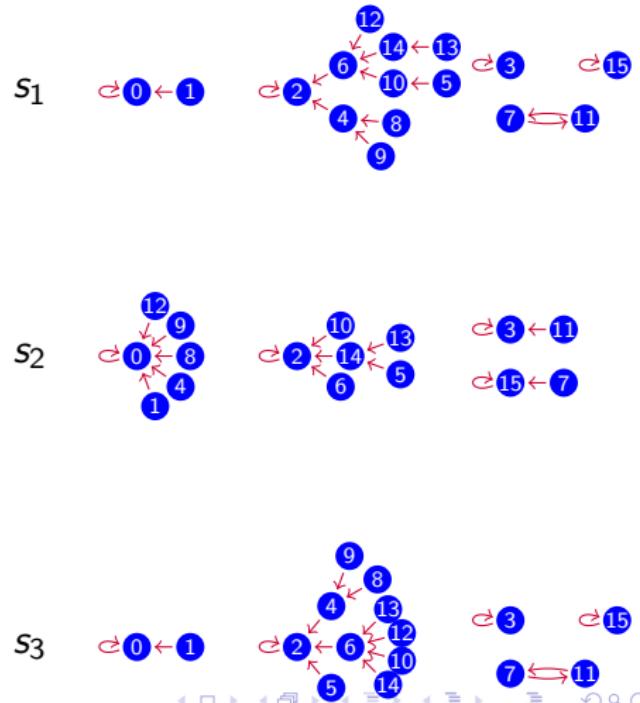
# Dynamics

$$s_1 = \{1, 2, 3, 4\}$$

$$s_2 = \{1\} \{2\} \{3\} \{4\}$$

$$s_3 = \{3, 4\} \{1, 2\}$$

	State	$s_1$	$s_2$	$s_3$
0	0000	0000	0000	0000
1	0001	0000	0000	0000
2	0010	0010	0010	0010
3	0011	0011	0011	0011
4	0100	0010	0000	0010
5	0101	1010	1110	0010
6	0110	0010	0010	0010
7	0111	1011	1111	1011
8	1000	0100	0000	0100
9	1001	0100	0000	0100
10	1010	0110	0010	0110
11	1011	0111	0011	0111
12	1100	0110	0000	0110
13	1101	1110	1110	0110
14	1110	0110	0010	0110
15	1111	1111	1111	1111

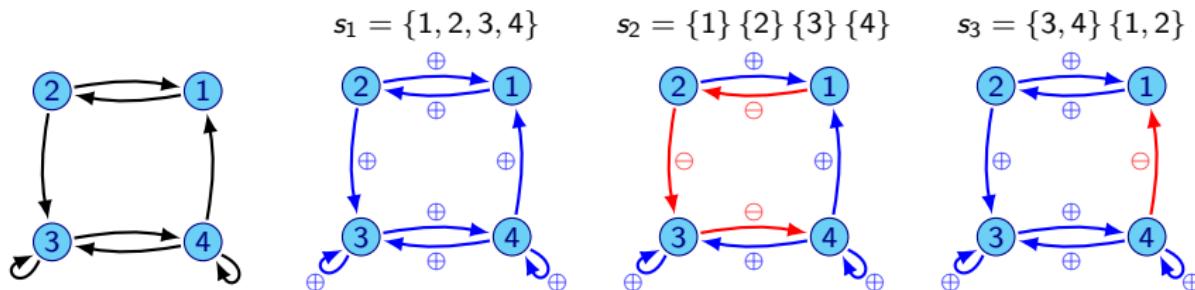


# Labeled digraph

A labeled digraph is a graph  $G$  with a label function  $lab$ ,  $(G, lab)$  such that:  $lab : A(G) \rightarrow \{\oplus, \ominus\}$

We say that a labeled digraph is an update digraph if there exists  $s : V(G) \rightarrow \{1, \dots, n\}$ , an update function such that:

$$\forall (u, v) \in A(G), lab(u, v) = \oplus \iff s(u) \geq s(v)$$

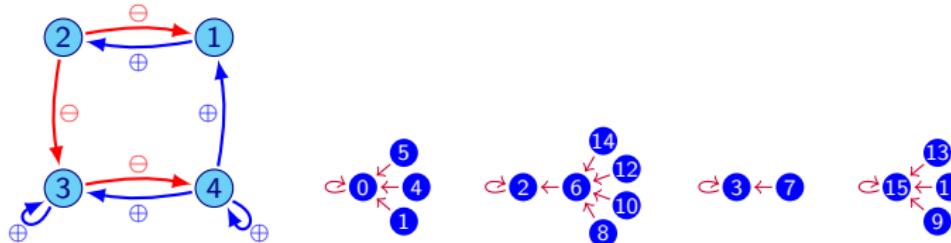


# Why are we interested in labeled digraphs?

Theorem. (Aracena et al (2008))

Given two Boolean networks  $N_1 = (F, s)$  and  $N_2 = (F, s')$  which differ only in the update schedule. If the labeled digraphs associated to them are equal, then both networks have the same dynamical behavior.

$$s = \{2\} \{1\} \{3\} \{4\} \text{ and } s' = \{2\} \{3\} \{1\} \{4\}$$



# Labels and Update digraphs

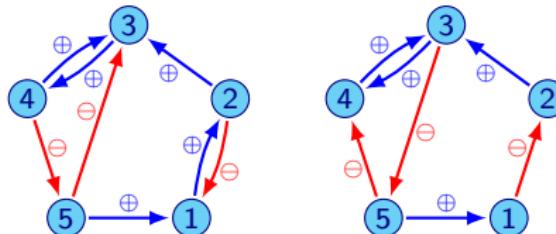
## Forbidden cycle

A forbidden cycle is a cycle with an arc labeled  $\ominus$  in the reverse graph.

## Theorem (Montalva (2012))

A labeled digraph is an update digraph if and only if there does not exist a forbidden cycle in its reverse digraph.

Labeled digraph    Reverse digraph



# Labels and Update digraphs

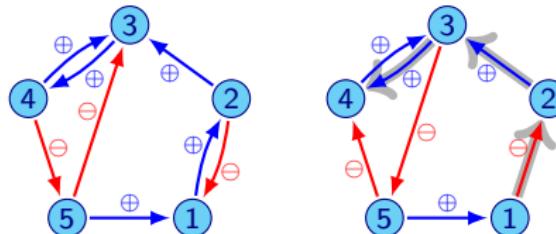
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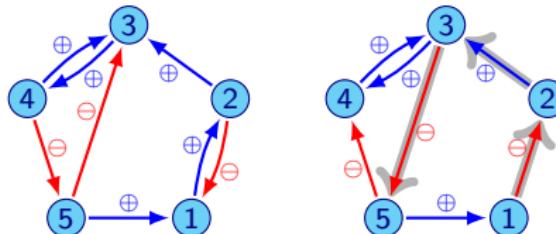
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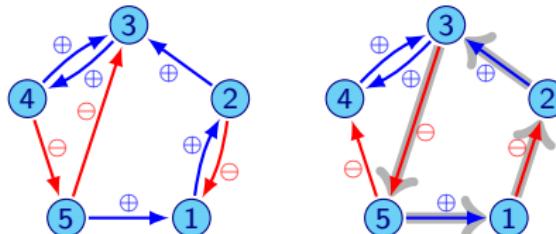
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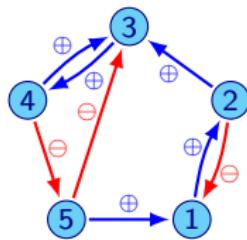
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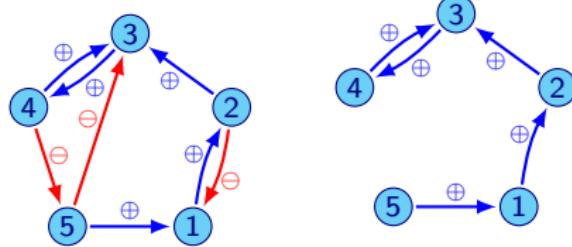
Labeled digraph    Reverse digraph



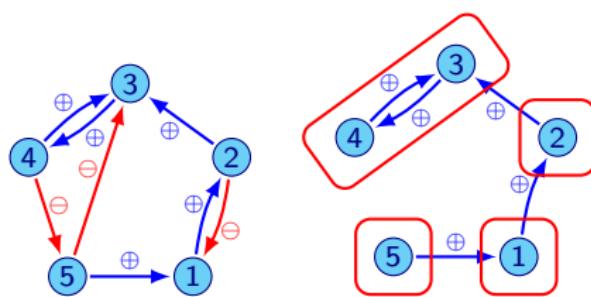
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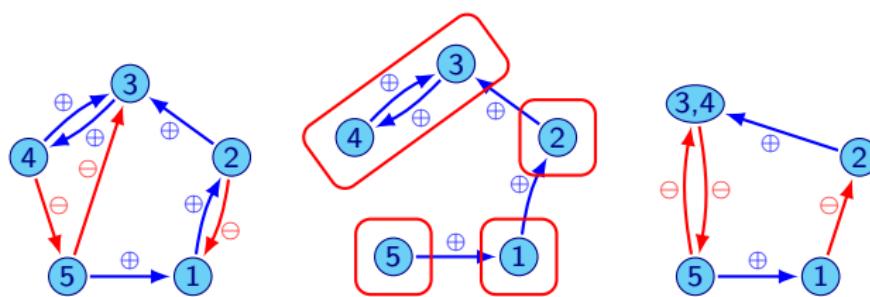
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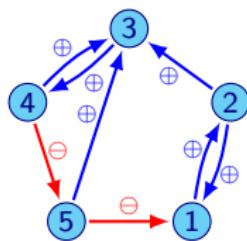
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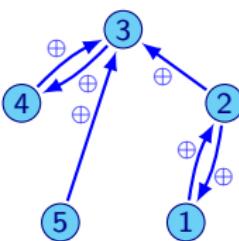
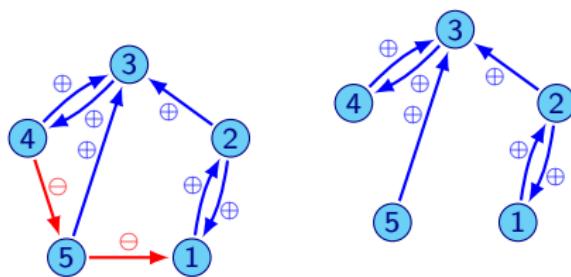
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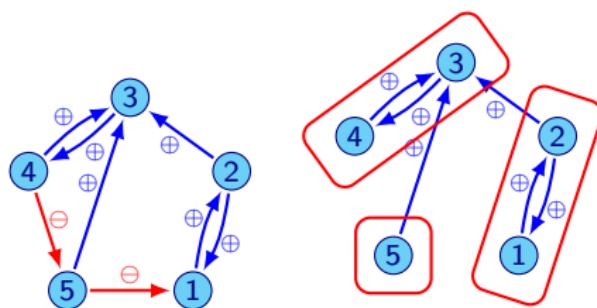
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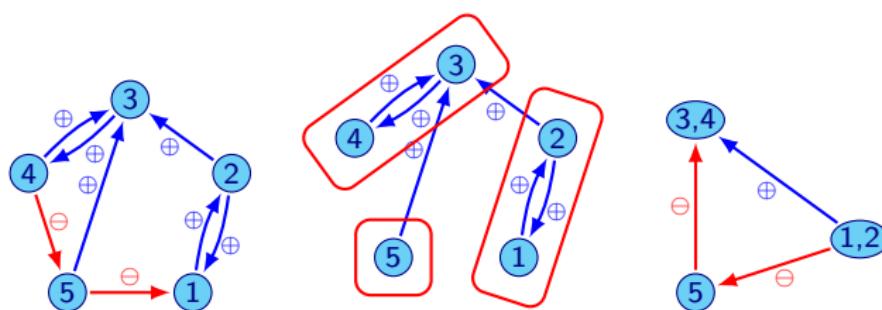
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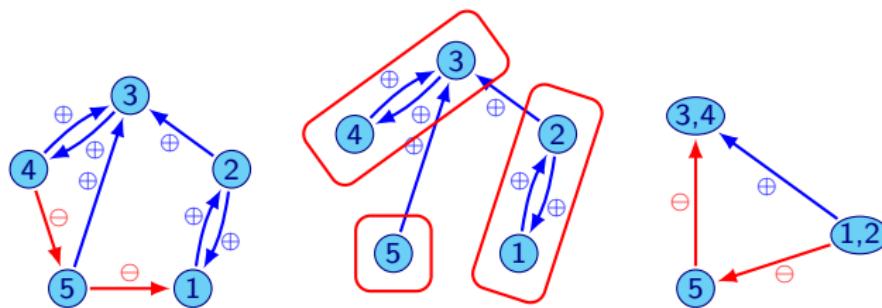
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$$s = \{3, 4\} \{5\} \{1, 2\}$$

# Update Digraph Extension Problem

## UDE

Given a labeled digraph  $(G, \text{lab})$ , find the set  $\mathcal{S}(G, \text{lab})$  of all fully labeled extensions  $\text{lab}'$  of  $\text{lab}$  such that  $(G, \text{lab}')$  is an update digraph.

# Complexity

## CUDE

Given  $(G, \text{lab})$  a labeled digraph, to determine the cardinality of the set  $\mathcal{S}(G, \text{lab})$ .

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## Theorem

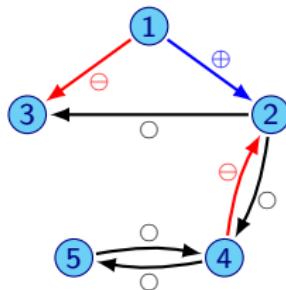
CUDE is  $\#P$ -complete

# Force

## Proposition

Given a labeled digraph  $(G, \text{lab})$  and an arc  $(i, j)$  with  $\text{lab}(i, j) = \circlearrowright$ :

- If there exists a reverse path from  $i$  to  $j$ , then the arc  $(i, j)$  must be labeled  $\oplus$
- If there exists a negative reverse path from  $j$  to  $i$ , then the arc  $(i, j)$  must be labeled  $\ominus$

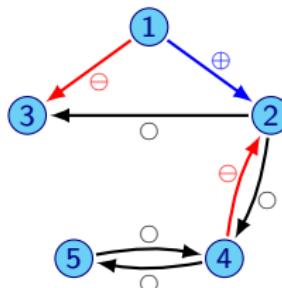


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Matrix  $M$

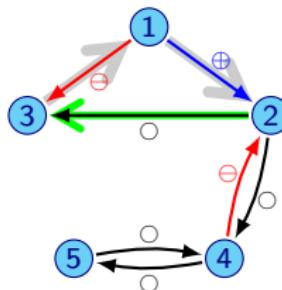
$V(G)$	1	2	3	4	5
1	$\infty$	1	$\infty$	-1	$\infty$
2	$\infty$	$\infty$	$\infty$	-1	$\infty$
3	-1	-1	$\infty$	-1	$\infty$
4	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$
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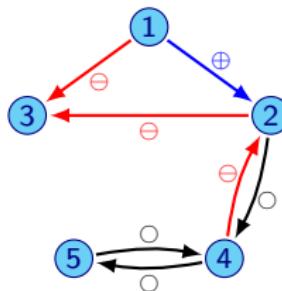
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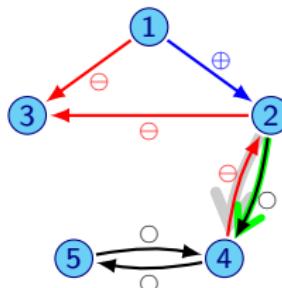
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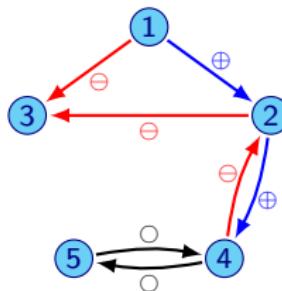
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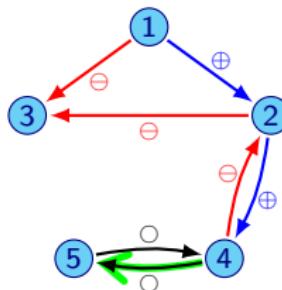
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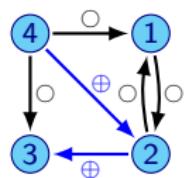
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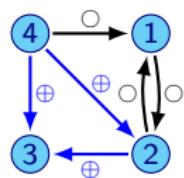
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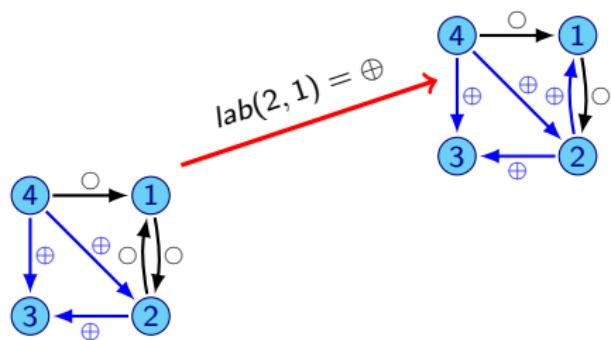
# Example of SimpleLabel



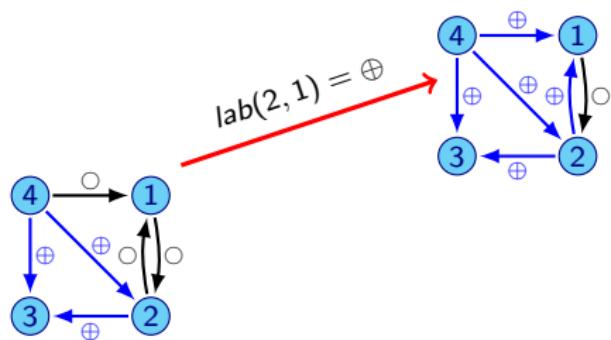
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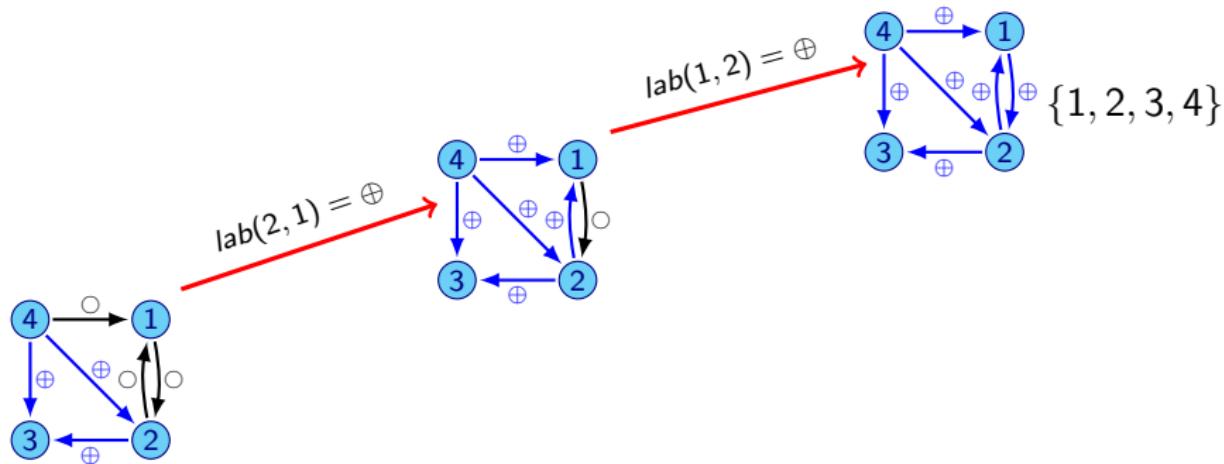
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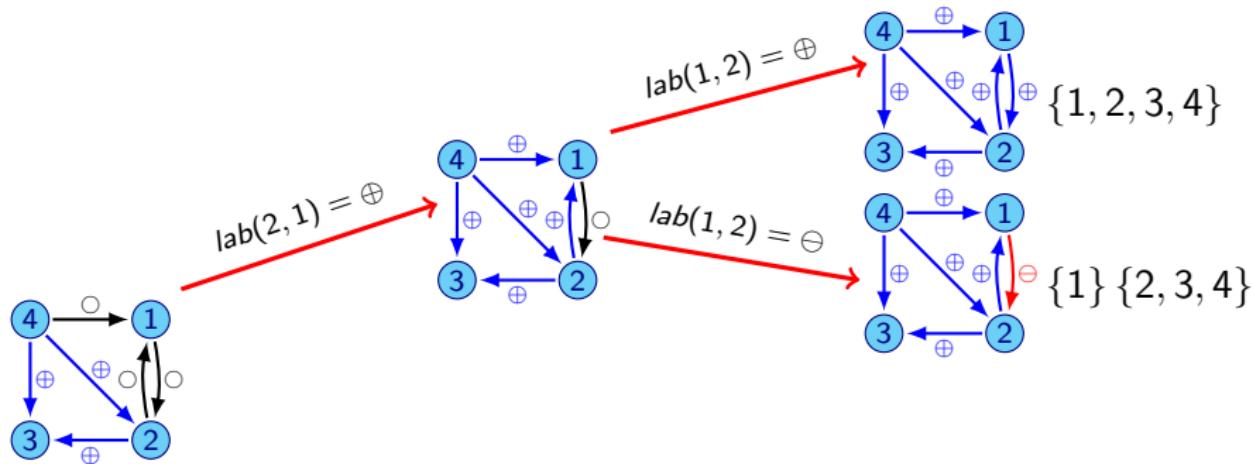
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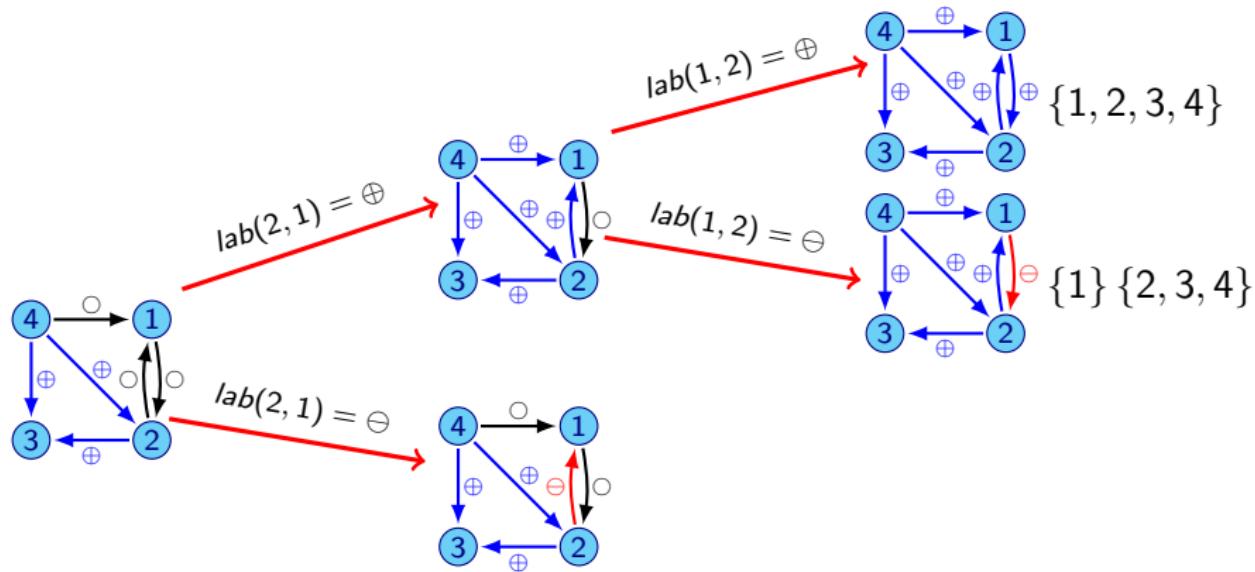
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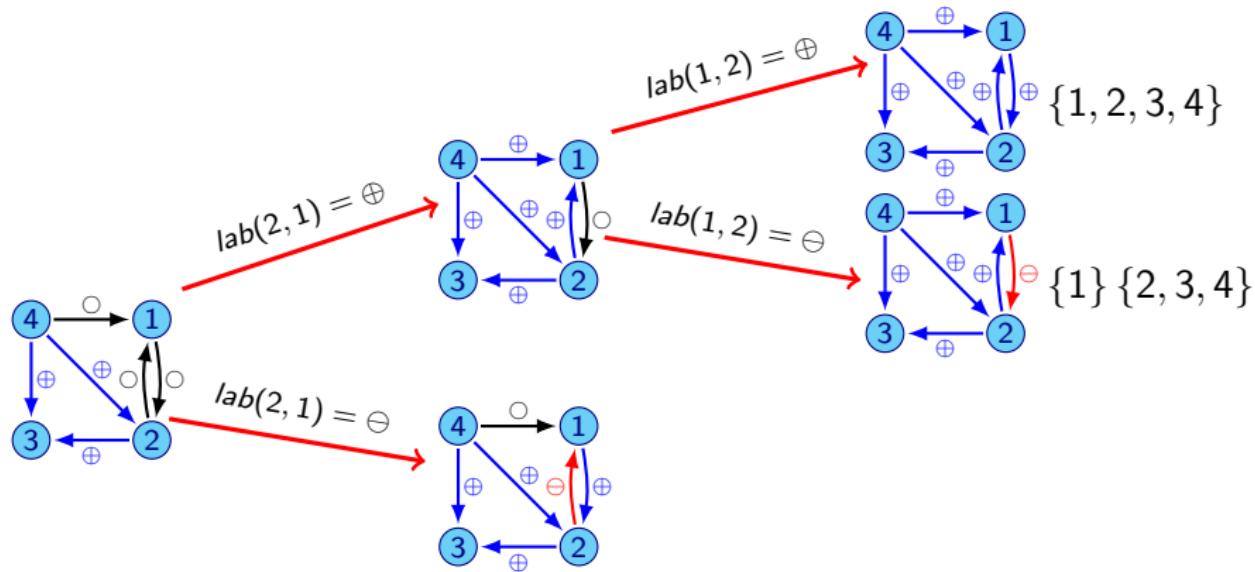
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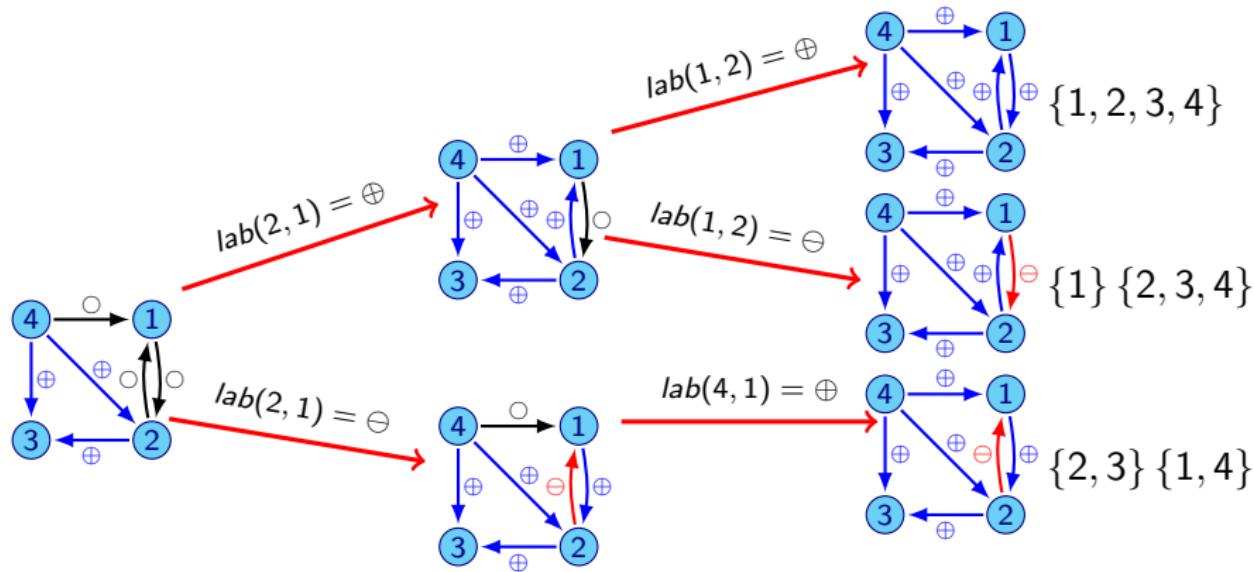
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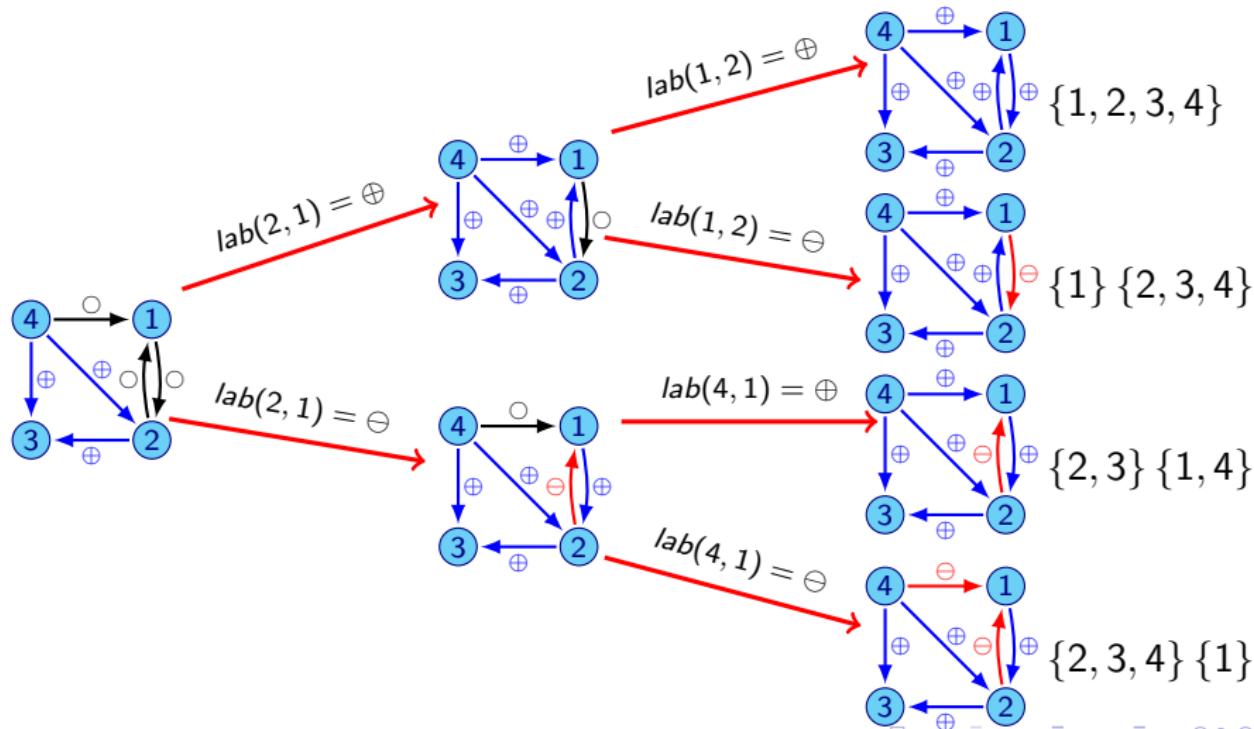
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# Reduce

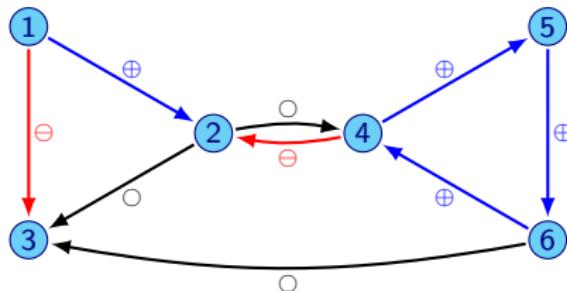
## Definition

Let  $(G, \text{lab})$  be an update digraph and  $\{G_1, \dots, G_k\}$  its positive strongly connected components. We define its reduced labeled digraph by

$R(G, \text{lab}) = (G_{rd} = (V_{rd}, A_{rd}), \text{lab}_{rd})$ , where:

- $V_{rd} = \{v_1, \dots, v_k\}$
- $A_{rd} = \{(v_i, v_j) | \exists (u, v) \in A(G) \cap (V(G_i) \times V(G_j))\}$

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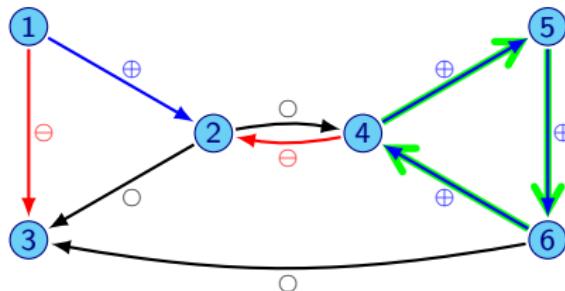
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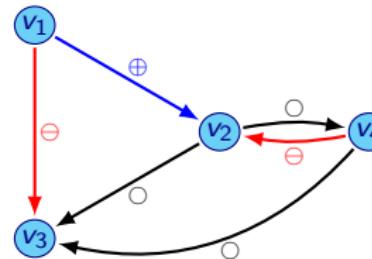
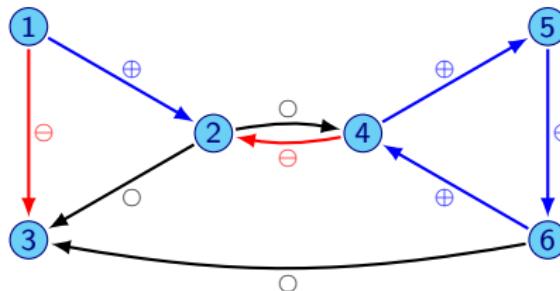
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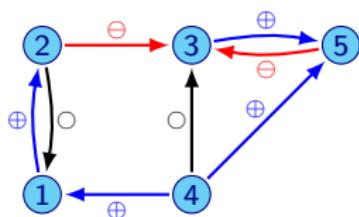
# Divide and conquer

## Divide by SCC

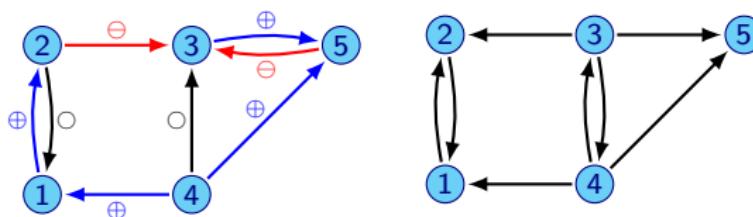
Let  $(G, \text{lab})$  an update digraph with SCC  $G_1, \dots, G_k$  (ordered) over its reverse extended digraph, then:

$$\mathcal{S}(G, \text{lab}) = \mathcal{S}\left(\tilde{G}_1, \text{lab}|_{A(G_1)}\right) \circ_n \cdots \circ_n \mathcal{S}\left(\tilde{G}_k, \text{lab}|_{A(G_k)}\right)$$

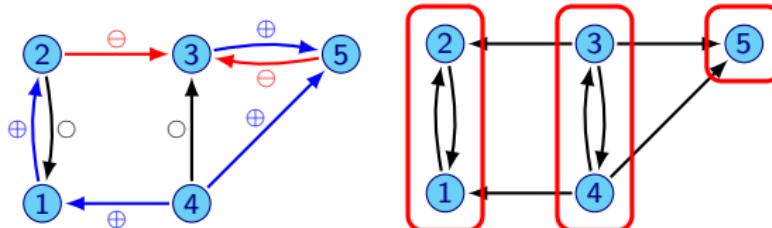
# Example: Division by SCC



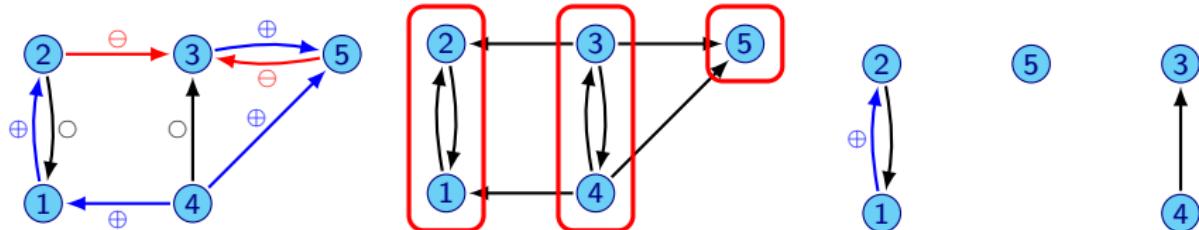
# Example: Division by SCC



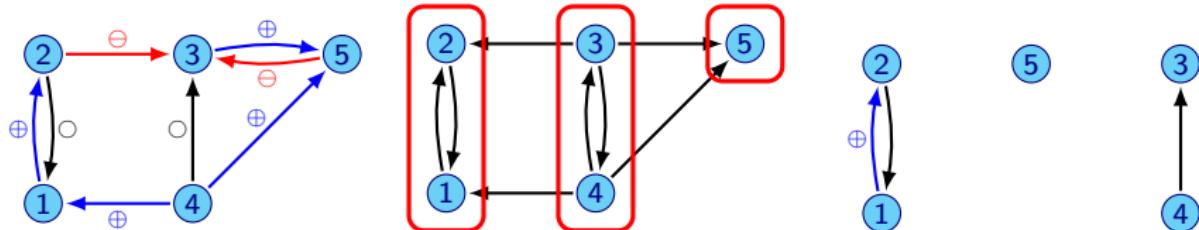
# Example: Division by SCC



# Example: Division by SCC



# Example: Division by SCC

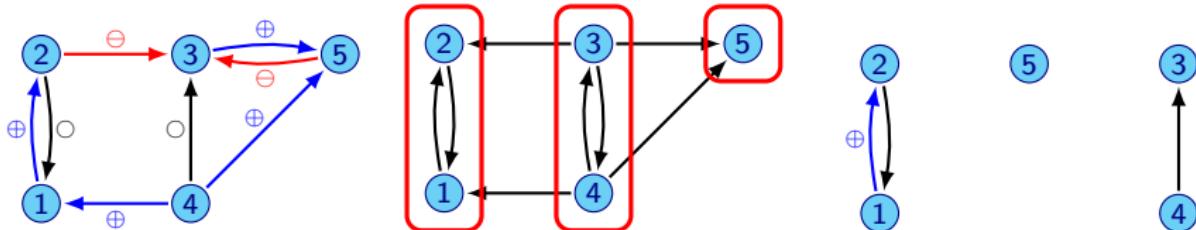


$$\mathcal{S}(G[1, 2], \text{lab}) = \{\{1, 2\}, \{2\} \{1\}\}$$

$$\mathcal{S}(G[5], \text{lab}) = \{\{5\}\}$$

$$\mathcal{S}(G[3, 4], \text{lab}) = \{\{3\} \{4\}, \{4\} \{3\}\}$$

# Example: Division by SCC



$$\mathcal{S}(G[1, 2], \text{lab}) = \{\{1, 2\}, \{2\} \{1\}\}$$

$$\mathcal{S}(G[5], \text{lab}) = \{\{5\}\}$$

$$\mathcal{S}(G[3, 4], \text{lab}) = \{\{3\} \{4\}, \{4\} \{3\}\}$$

Then,

$$\begin{aligned} \mathcal{S}(G, \text{lab}) &= \mathcal{S}(G[1, 2], \text{lab}) \circ_n \mathcal{S}(G[5], \text{lab}) \circ_n \mathcal{S}(G[3, 4], \text{lab}) \\ &= \{\{1, 2\} \{5\}, \{1\} \{2\} \{5\}\} \circ_n \{\{3\} \{4\}, \{4\} \{3\}\} \\ &= \left\{ \{1, 2\} \{5\} \{3\} \{4\}, \{1\} \{2\} \{5\} \{3\} \{4\}, \right. \\ &\quad \left. \{1, 2\} \{5\} \{4\} \{3\}, \{1\} \{2\} \{5\} \{4\} \{3\} \right\} \end{aligned}$$

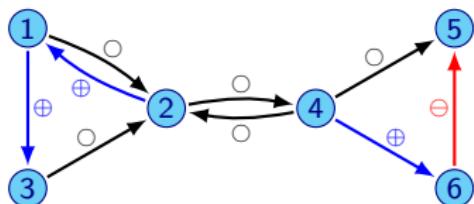
# Divide and conquer

## Divide by Bridges

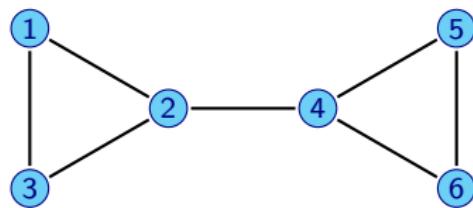
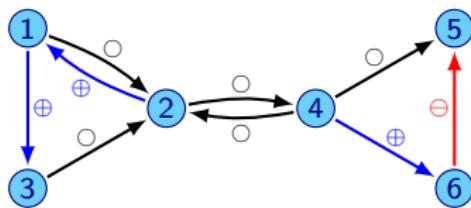
Let  $(G, \text{lab})$  a connected digraph,  $G_U$  the underlying digraph of  $G$  and  $uv \in E(G_U)$  a bridge that divide  $G$  in  $G_1$  and  $G_2$ , then

$$\mathcal{S}(G, \text{lab}) = \mathcal{S}(G_1, \text{lab}|_{A(G_1)}) \circ_{\{u,v\}} \mathcal{S}(G_2, \text{lab}|_{A(G_2)})$$

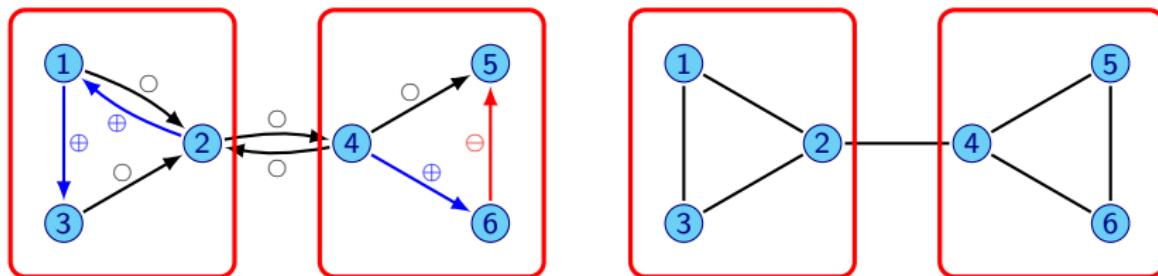
# Example: Division by Bridges



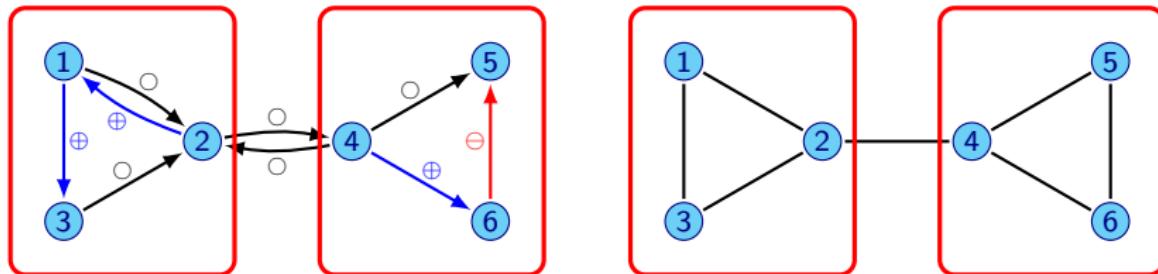
# Example: Division by Bridges



# Example: Division by Bridges



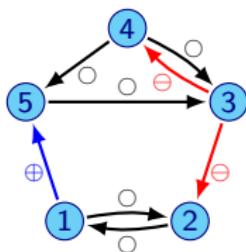
# Example: Division by Bridges



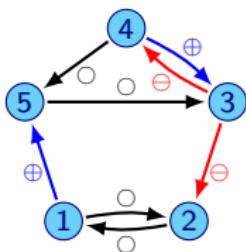
$$S_1 \circ_{nrm_{2,4}} S_2 = \{\{1, 2, 3\}, \{3\} \{1, 2\}, \{3\} \{1\} \{2\}\} \circ_{nrm_{2,4}} \{\{6\} \{5\} \{4\}, \{6\} \{4\} \{5\}\}$$

$\{1, 2, 3\} \{6\} \{5\} \{4\}$	$\{6\} \{5\} \{4\} \{1, 2, 3\}$	$\{6\} \{5\} \{4, 1, 2, 3\}$
$\{3\} \{1, 2\} \{6\} \{5\} \{4\}$	$\{6\} \{5\} \{4\} \{3\} \{1, 2\}$	$\{3\} \{6\} \{5\} \{4, 1, 2\}$
$\{3\} \{1\} \{2\} \{6\} \{5\} \{4\}$	$\{6\} \{5\} \{4\} \{3\} \{1\} \{2\}$	$\{3\} \{1\} \{6\} \{5\} \{4, 2\}$
$\{1, 2, 3\} \{6\} \{4\} \{5\}$	$\{6\} \{4\} \{5\} \{1, 2, 3\}$	$\{6\} \{4, 1, 2, 3\} \{5\}$
$\{3\} \{1, 2\} \{6\} \{4\} \{5\}$	$\{6\} \{4\} \{5\} \{3\} \{1, 2\}$	$\{3\} \{6\} \{4, 1, 2\} \{5\}$
$\{3\} \{1\} \{2\} \{6\} \{4\} \{5\}$	$\{6\} \{4\} \{5\} \{3\} \{1\} \{2\}$	$\{3\} \{1\} \{6\} \{4, 2\} \{5\}$

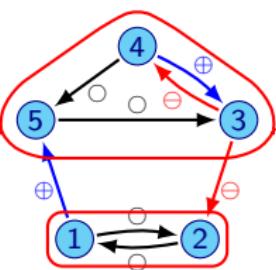
# Example: UpdateLabel



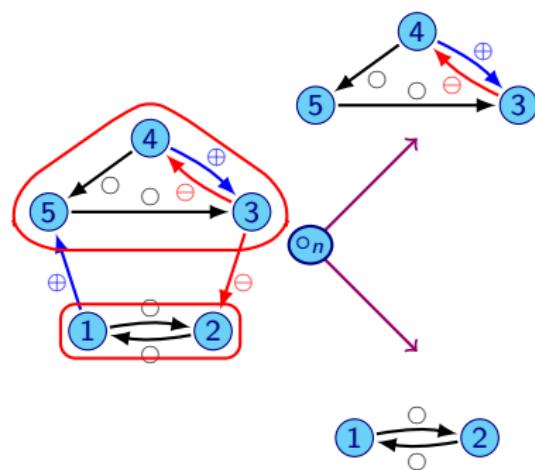
# Example: UpdateLabel



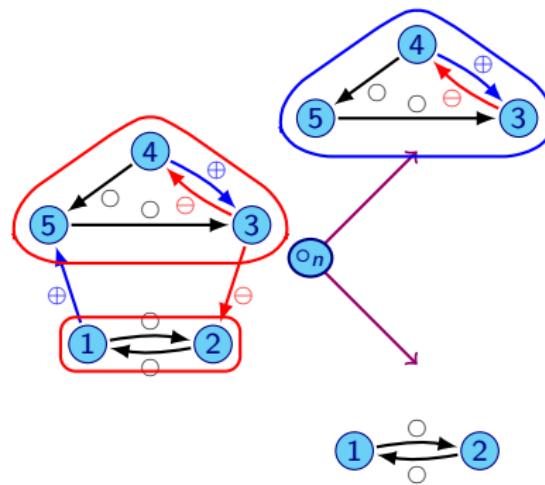
# Example: UpdateLabel



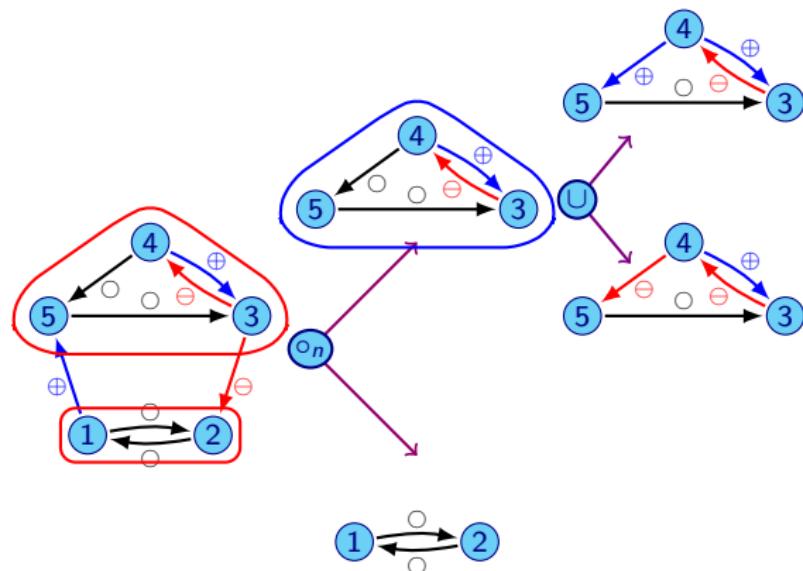
# Example: UpdateLabel



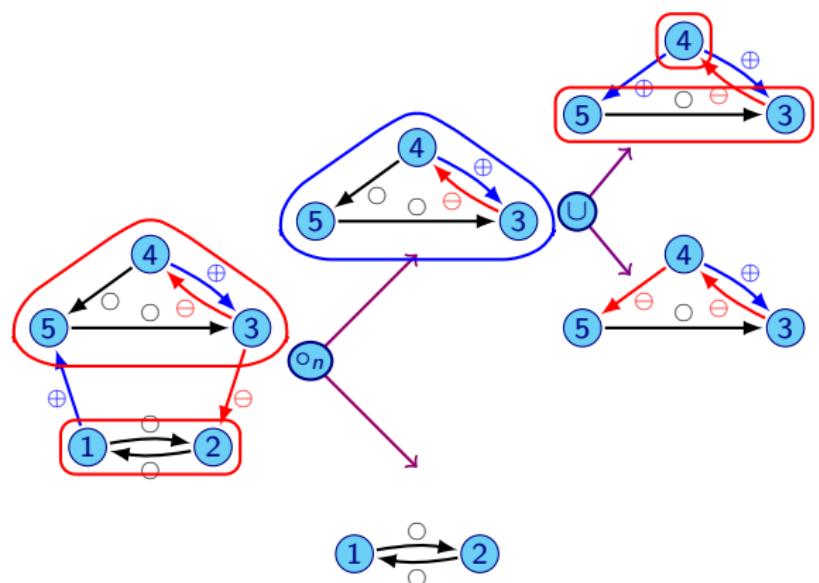
# Example: UpdateLabel



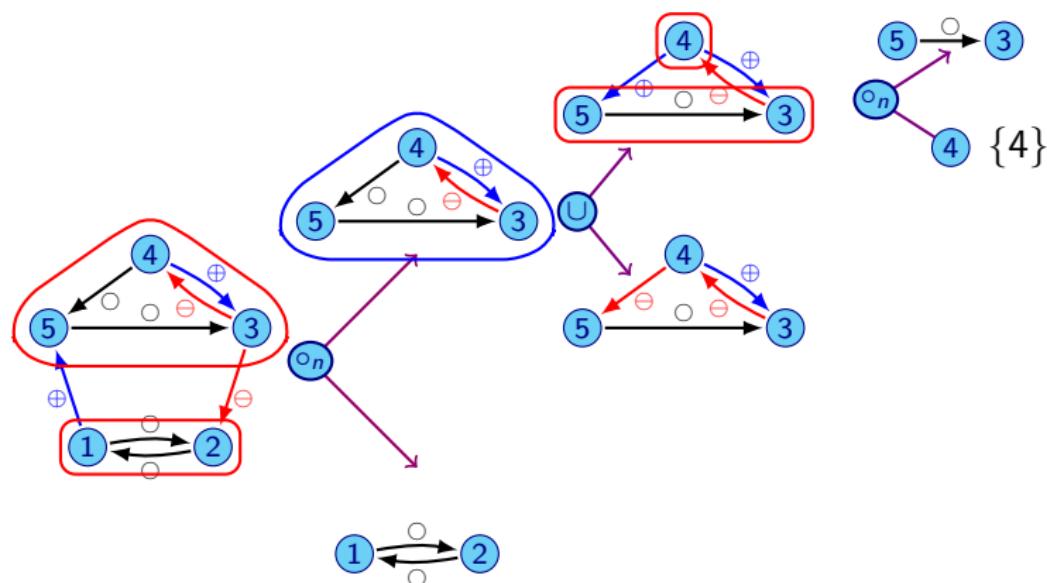
# Example: UpdateLabel



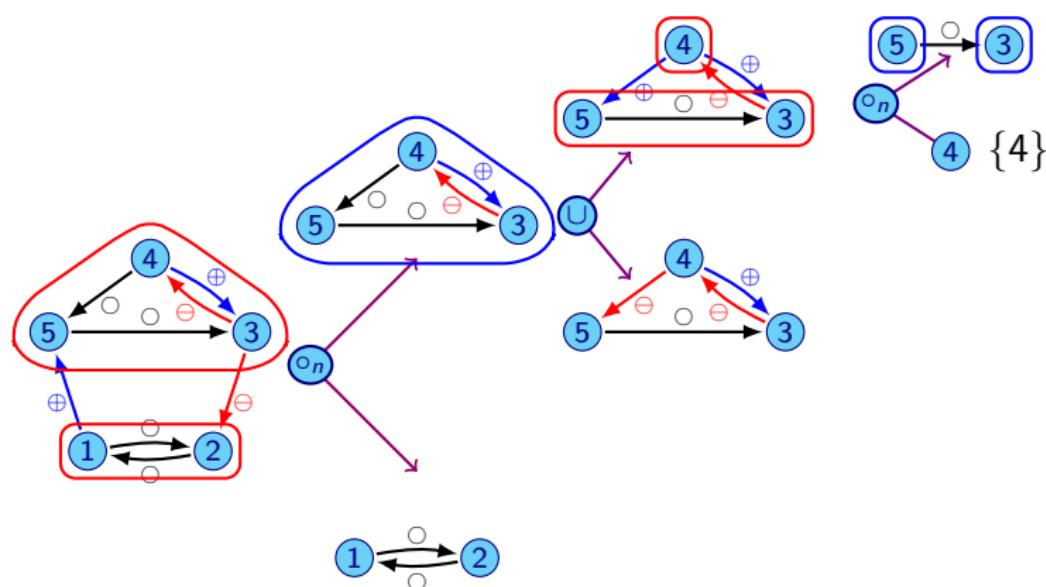
# Example: UpdateLabel



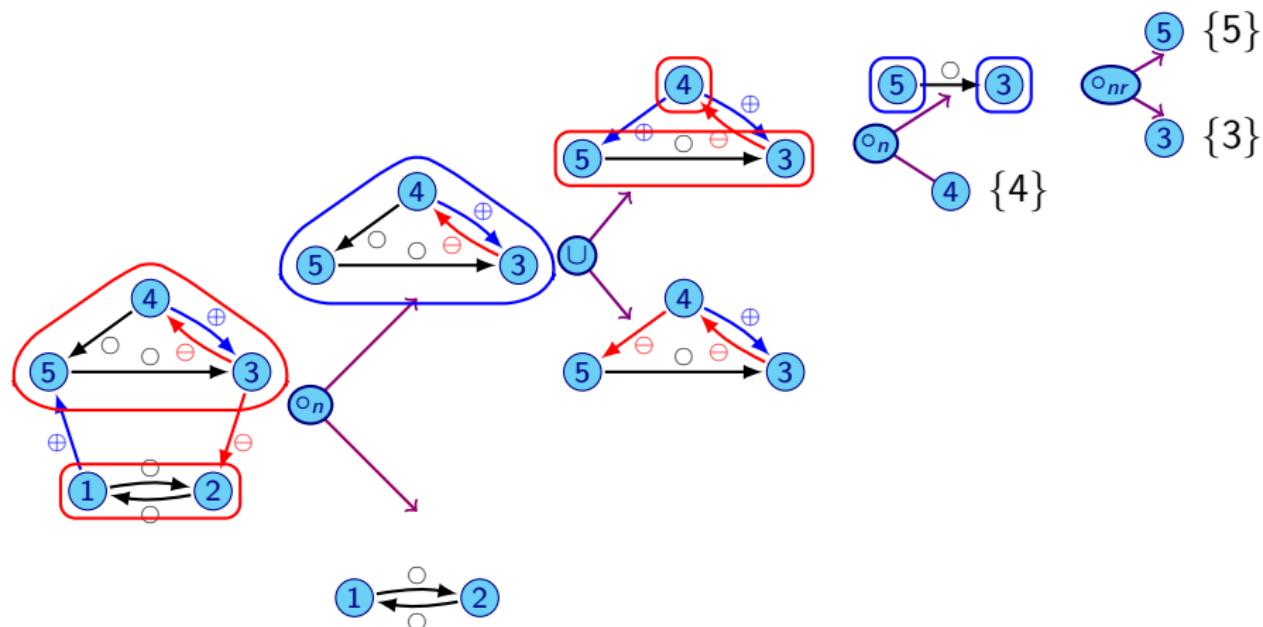
# Example: UpdateLabel



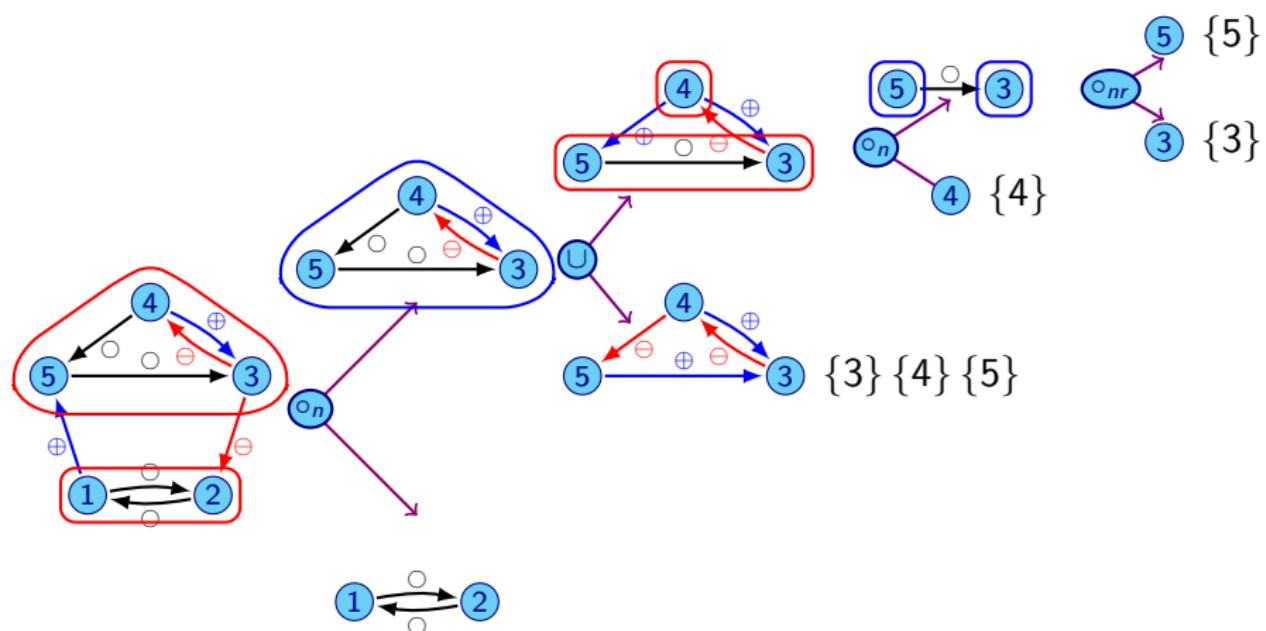
## Example: UpdateLabel



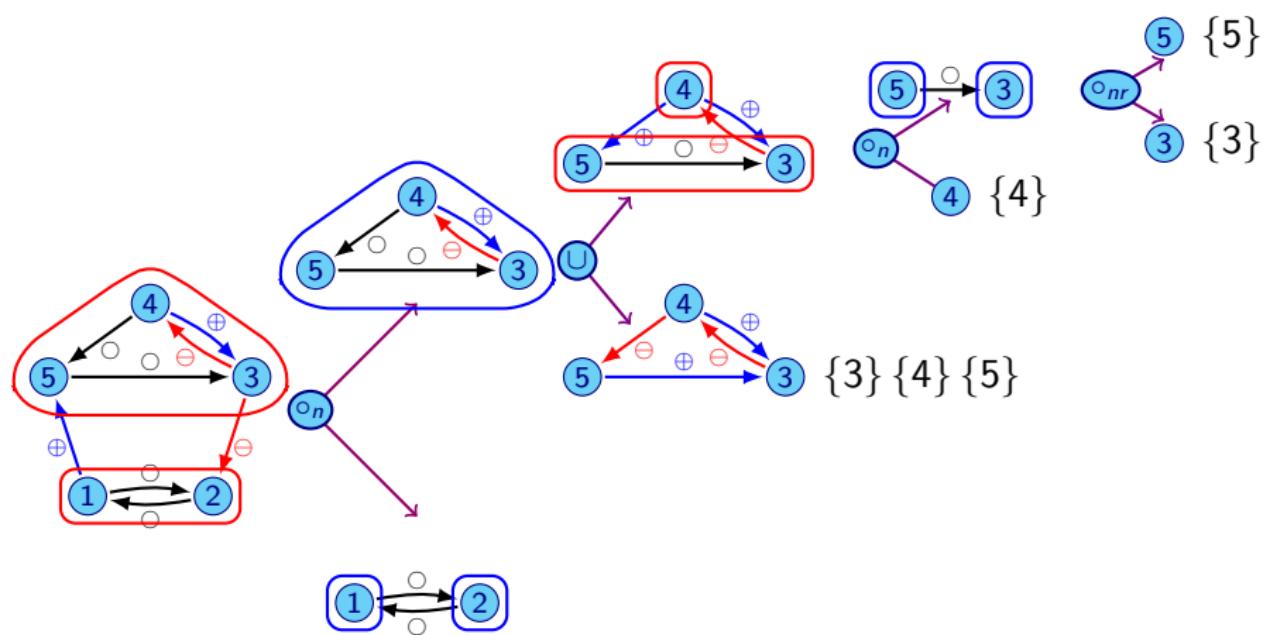
## Example: UpdateLabel



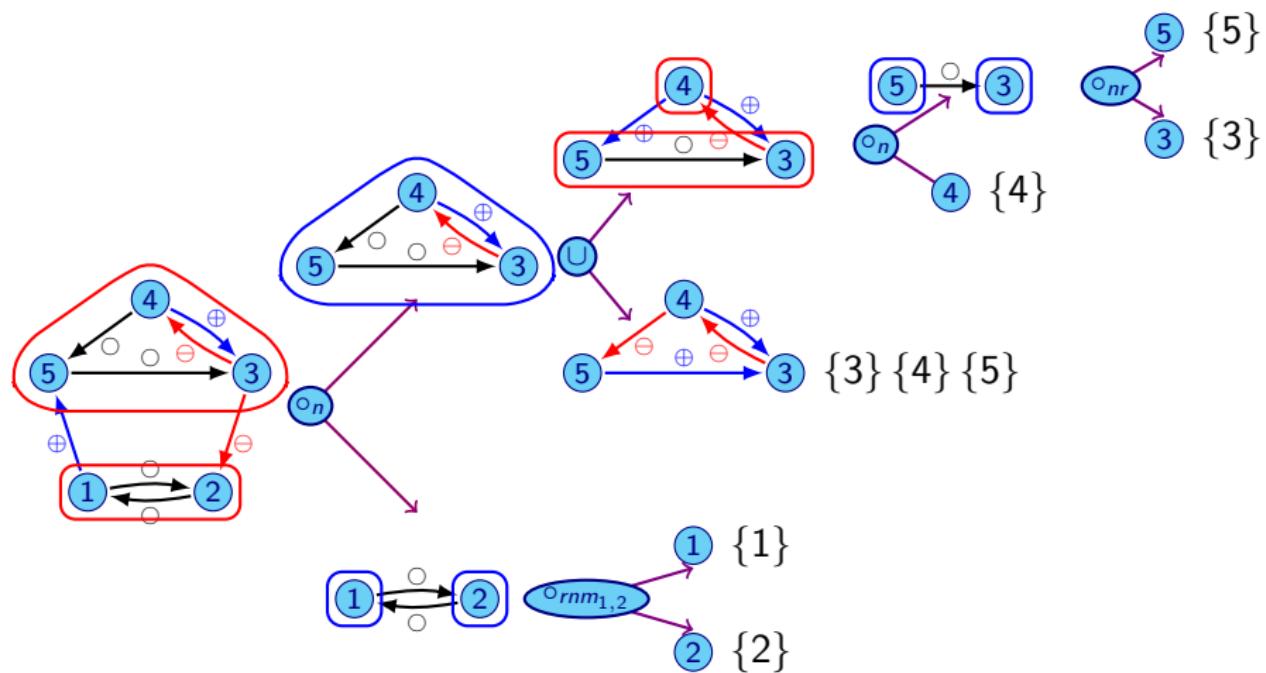
# Example: UpdateLabel



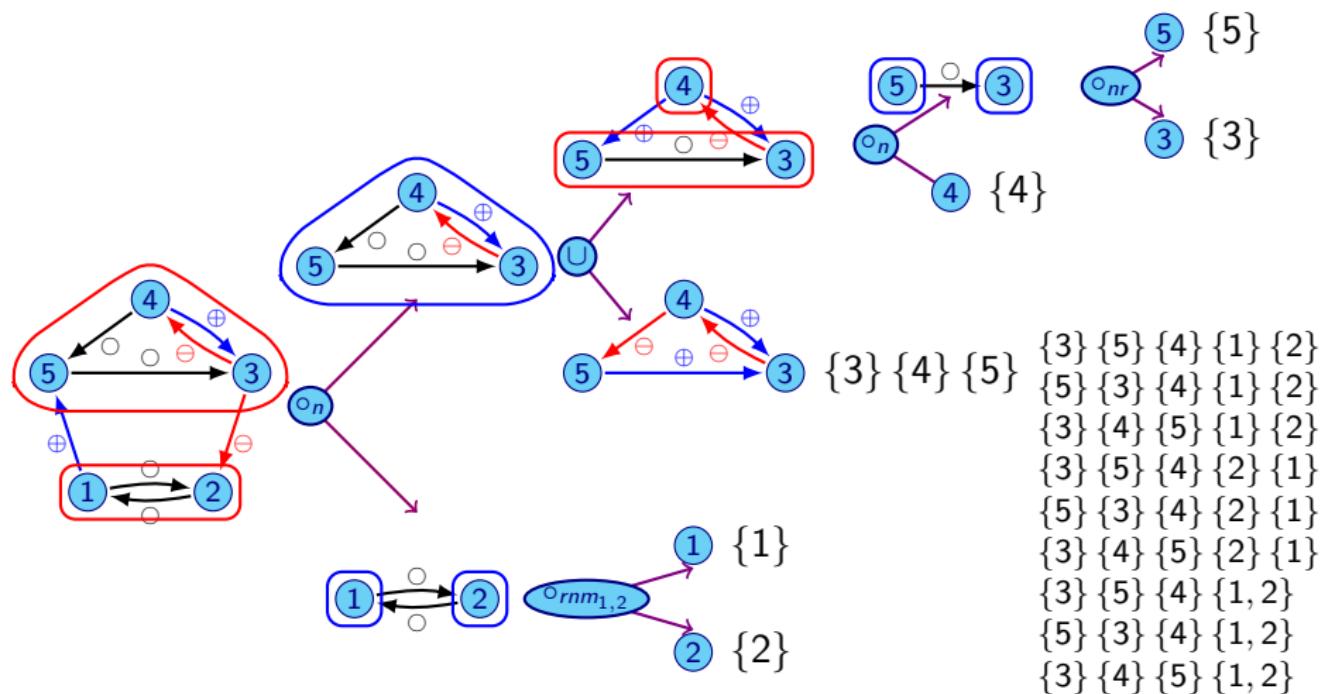
# Example: UpdateLabel



# Example: UpdateLabel



# Example: UpdateLabel



# Results

Graph	Nodes	Arcs	$ \mathcal{S}(G, \text{lab}) $	SimpleLabel	Label
$K_3$	3	6	13	0.02	0.02
$K_5$	5	20	541	0.18	0.13
$K_7$	7	42	47 293	140.76	0.79
$K_8$	8	56	545 835	> 22 200.00	1.90
$P_{10}$	10	18	19 683	4.43	0.02
$P_{12}$	12	22	177 147	313.68	0.02
$P_{50}$	50	98	$3^{49}$	—	0.09
$P_{200}$	200	398	$3^{199}$	—	0.17