

14. Algoritmos de Ordenamiento

Input: una secuencia de n elementos $(a_1, \dots, a_n)$

Output: una permutación (reordenación) $(aa_1, \dots, aa_2)$ de la secuencia de entrada tal que $aa_1 \leq \dots \leq aa_2$

Tipos de algoritmos de ordenamiento:

- Ordenamiento por Selección

```
Procedure Selection_Sort ( $A$ )
1   for  $i = 1$  to  $length[A]$  do [
2        $min := MAX\_INTEGER$ 
3       for  $j = i$  to  $length[A]$  do [
4           if  $A[j] < min$  then [
5                $imin := j$ 
6                $min := A[j]$ ]
7        $temp := A[i]$ 
8        $A[i] := A[imin]$ 
9        $A[imin] := temp$ 
end Selection_Sort
```

- Ordenamiento por Inserción

```
Procedure Insertion_Sort ( $A$ )
1   for  $j = 2$  to  $length[A]$  do [
2        $i := j - 1$ 
3        $key := A[j]$ 
4       while  $i > 0$  and  $A[i] > key$  do [
5            $A[i + 1] := A[i]$ 
6            $i := i - 1$ 
7        $A[i + 1] := key$ ]
end Insertion_Sort
```

- Ordenamiento por Burbuja

```
Procedure Bubble_Sort ( $A$ )
1   for  $i = 1$  to  $length[A]$  do [
2       for  $j = 1$  to  $length[A] - i$  do [
3           if  $A[j] > A[j + 1]$  then [
4                $temp := A[j]$ 
5                $A[j] := A[j + 1]$ 
6                $A[j + 1] := temp$ ]
end Bubble_Sort
```

- Ordenamiento por Mezcla

Procedure Merge_Sort (A, p, r)

```

1      if  $p < r$  then [
2           $q := \lfloor (p + r)/2 \rfloor$ 
3          Merge_Sort( $A, p, q$ )
4          Merge_Sort( $A, q + 1, r$ )
5          Merge( $A, p, q, r$ )]
end Merge_Sort

```

Procedure Merge(A, p, q, r)

```

1       $n_1 := q - p + 1$ 
2       $n_2 := r - q$ 
3      for  $i = 1$  to  $n_1$  do  $L[i] := A[p + i - 1]$ 
4      for  $i = 1$  to  $n_2$  do  $R[i] := A[q + i]$ 
5       $L[n_1 + 1] = \infty$ 
6       $R[n_2 + 1] = \infty$ 
7       $i := 1$ 
8       $j := 1$ 
9      for  $k = p$  to  $r$  do
10         if  $L[i] \leq R[j]$  then[
11              $A[k] := L[i]$ 
12              $i := i + 1$ ]
13         else
14              $A[k] := R[j]$ 
15              $j := j + 1$ ]
end Merge

```

- Ordenamiento Rápido

Procedure QuickSort (A, p, r)

```

1      if  $p < r$  then [
2           $q := \text{Partition}(A, p, r)$ 
3          QuickSort( $A, p, q$ )
4          QuickSort( $A, q + 1, r$ ) ]
end QuickSort

```

```

Procedure Partition( $A, p, r$ )
1       $x := A[p]$ 
2       $i := p - 1$ 
3       $j := r + 1$ 
4      while true do [
5          repeat  $j := j - 1$ 
6              until  $A[j] \leq x$ 
7          repeat  $i := i + 1$ 
8              until  $A[i] \geq x$ 
9          if  $i < j$  then[
10              $temp := A[j]$ 
11              $A[j] := A[i]$ 
12              $A[i] := temp$ 
13         else return  $j$  ]
end Partition

```

■ Ordenamiento Heap

```

Procedure HeapSort( $A, p, r$ )
1      Build_Heap( $A$ )
2      for  $i = length(A)$  downto 2 do [
3           $temp := A[1]$ 
4           $A[1] = A[i]$ 
5           $A[i] = temp$ 
6           $heap\_size(A) := heap\_size(A) - 1$ 
7          Heapify( $A, 1$ ) ]
end HeapSort

```

```

Procedure Build_Heap( $A$ )
     $heap\_size := length(A)$ 
    for  $i = (length(A)/2)$  downto 1 do
        HEAPIFY( $A, i$ )
end Build_Heap

```

```

Procedure Heapify( $A, i$ )
     $l := Left(i)$ 
     $r := Right(i)$ 
    if  $l \leq heap\_size(A)$  and  $A[l] > A[i]$  then
         $largest := l$ 
    else
         $largest := i$ 
    if  $r \leq heap\_size(A)$  and  $A[r] > A[largest]$  then
         $largest := r$ 
    if  $largest \neq i$  then [
        exchange( $A[i], A[largest]$ )
        Heapify( $A, largest$ ) ]
end Heapify

```

■ Ordenamiento Rápido_2

```
Procedure QuickSort_2 ( $A, p, r$ )  
1      if  $p < r$  then [  
2           $q := Partition(A, p, r)$   
3           $QuickSort(A, p, q - 1)$   
4           $QuickSort(A, q + 1, r)$  ]  
end QuickSort_2
```

```
Procedure Partition_2( $A, p, r$ )  
1           $x := A[r]$   
2           $i := p - 1$   
3          for  $j := p$  to  $r - 1$  do [  
4              if  $A[j] \leq x$  then [  
5                   $i := i + 1$   
6                  exchange  $A[i] \leftrightarrow A[j]$ ]  
7              exchange  $A[i + 1] \leftrightarrow A[r]$   
8              return  $i + 1$   
end Partition_2
```